Distributed Middleware

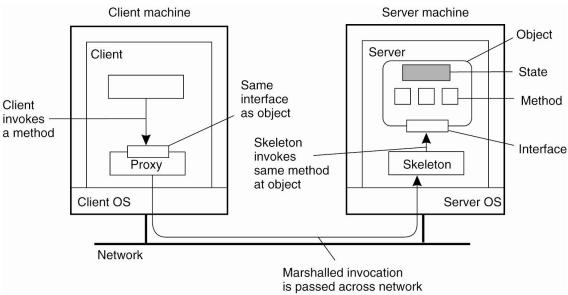
- Distributed objects
- DCOM
- CORBA
- EJBs
- Jini



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Distributed Objects



• Figure 10-1. Common organization of a remote object with client-side proxy.



Distributed Objects vs. RPC

RPC: Remote Procedure Call

- Provides argument marshalling / unmarshalling
- Server handles invocation

Distributed Objects

- Remote methods on remote objects
- RPC + distributed object references

Distributed object operation:

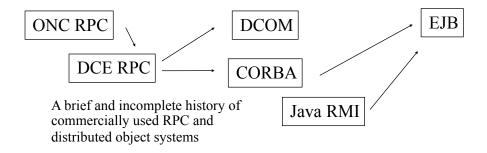
- Server side: create object, register it (register with what?) (always in this order?)
- Client side: get object reference (from where?), invoke method



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Distributed Objects through History

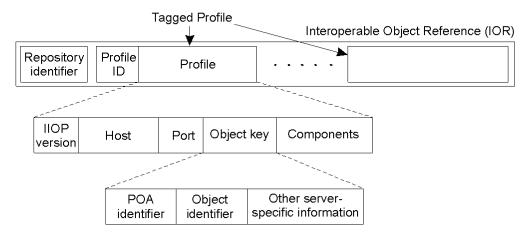




Computer Science

Naming: Object References

CORBA object reference



• Interoperable object reference: language-independent techniques for referring to objects



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Example: Enterprise Java Beans

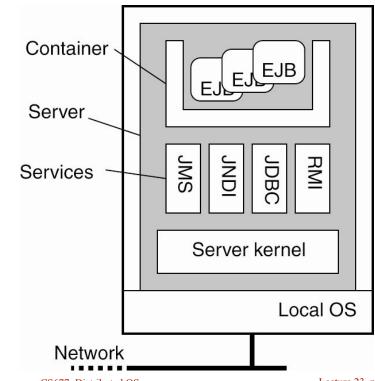


 Figure 10-2. General architecture of an EJB server.



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Parts of an EJB

- Home interface:
 - Object creation, deletion
 - Location of persistent objects (entity beans)
 - Object identifier is class-managed
- Remote interface
 - "business logic"
 - i.e. the object itself
- Terminology differences
 - Client/server -> web applications



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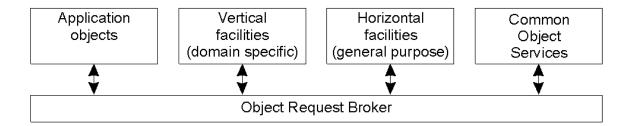
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Four Types of EJBs

- Stateless session beans
- Stateful session beans
- Entity beans
- Message-driven beans



CORBA Overview



- Object request broker (ORB)
 - Core of the middleware platform
 - Handles communication between objects and clients
 - Handles distribution and heterogeneity issues
 - May be implemented as libraries
- Facilities: composition of CORBA services



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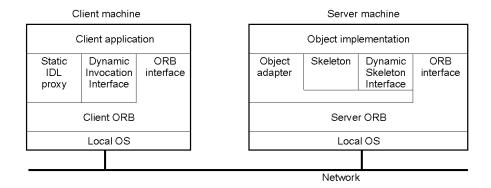
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Corba Services

Service	Description	
Collection	Facilities for grouping objects into lists, queue, sets, etc.	
Query	Facilities for querying collections of objects in a declarative manner	
Concurrency	Facilities to allow concurrent access to shared objects	
Transaction	Flat and nested transactions on method calls over multiple objects	
Event	Facilities for asynchronous communication through events	
Notification	Advanced facilities for event-based asynchronous communication	
Externalization	Facilities for marshaling and unmarshaling of objects	
Life cycle	Facilities for creation, deletion, copying, and moving of objects	
Licensing	Facilities for attaching a license to an object	
Naming	Facilities for systemwide name of objects	
Property	Facilities for associating (attribute, value) pairs with objects	
Trading	Facilities to publish and find the services on object has to offer	
Persistence	Facilities for persistently storing objects	
Relationship	Facilities for expressing relationships between objects	
Security	Mechanisms for secure channels, authorization, and auditing	
Time	Provides the current time within specified error margins	



Object Model



- Objects & services specified using an Interface Definition language (IDL)
 - Used to specify interface of objects and/or services
- ORB: run-time system that handles object-client communication
- Dynamic invocation interface: allows object invocation at run-time
 - Generic *invoke* operation: takes object reference as input
 - Interface repository stores all interface definitions



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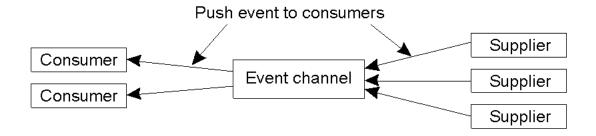
Object Invocation Models

Request type	Failure semantics	Description
Synchronous	At-most-once	Caller blocks until a response is returned or an exception is raised
One-way	Best effort delivery	Caller continues immediately without waiting for any response from the server
Deferred synchronous	At-most-once	Caller continues immediately and can later block until response is delivered

- Invocation models supported in CORBA.
 - Original model was RMI/RPC-like
 - Current CORBA versions support additional semantics



Event and Notification Services (1)



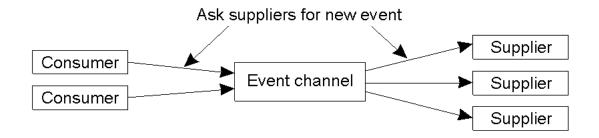
• The logical organization of suppliers and consumers of events, following the push-style model. (**PUB-SUB model**)



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Event and Notification Services (2)

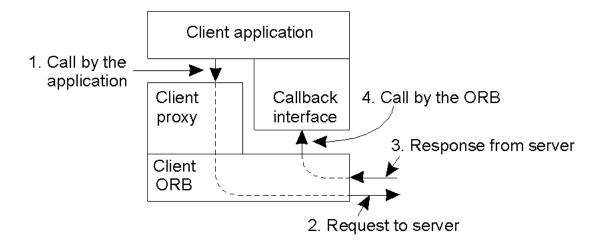


The pull-style model for event delivery in CORBA.



Messaging: Async. Method Invocation

CORBA's callback model for asynchronous method invocation.

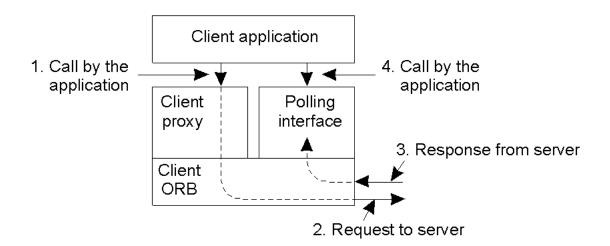




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Messaging (2)

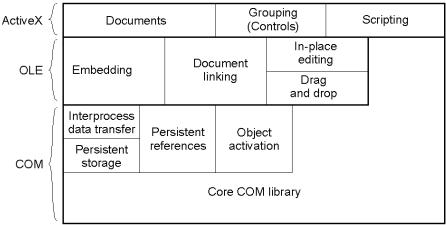


CORBA'S polling model for asynchronous method invocation.



DCOM

- Distributed Component Object Model
 - Microsoft's object model (middleware)
 - Now evolved into .NET





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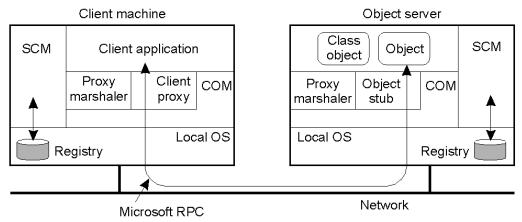
DCOM: History

- Successor to COM
 - Developed to support compound documents
 - Word document with excel spreadsheets and images
- Object linking and embedding (OLE)
 - Initial version: message passing to pass information between parts
 - Soon replaced by a more flexible layer: COM
- ActiveX: OLE plus new features
 - No good consensus on what exactly does ActiveX contain
 - Loosely: groups capabilities within applications to support scripting, grouping of objects.
- DCOM: all of the above, but across machines



Type Library and Registry

- The overall architecture of DCOM.
 - Type library == CORBA interface repository
 - Service control manager == CORBA implementation repository





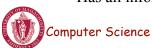
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Monikers: Persistent Objects

Step	Performer	Description	
1	Client	Calls BindMoniker at moniker	
2	Moniker	Looks up associated CLSID and instructs SCM to create object	
3	SCM	Loads class object	
4	Class object	Creates object and returns interface pointer to moniker	
5	Moniker	Instructs object to load previously stored state	
6	Object	Loads its state from file	
7	Moniker	Returns interface pointer of object to client	

- By default, DCOM objects are transient
- Persistent objects implemented using monikers (reference stored on disk)
 - Has all information to recreate the object at a later time



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Monikers (2)

Moniker type	Description	
File moniker	Reference to an object constructed from a file	
URL moniker	Reference to an object constructed from a URL	
Class moniker	Reference to a class object	
Composite moniker	Reference to a composition of monikers	
Item moniker	Reference to a moniker in a composition	
Pointer moniker	Reference to an object in a remote process	

• DCOM-defined moniker types.



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Distributed Coordination

- Motivation
 - Next generation of systems will be inherently distributed
 - Main problem: techniques to coordinate various components
 - Emphasis on coordination of activities between components



Introduction to Coordination Models

- Key idea: separation of computation from coordination
- A taxonomy of coordination models
 - Direct coordination
 - Mailbox coordination
 - Meeting-oriented coordination (publish/subscribe)
 - Generative (shared tuple space)

	Temporal	
	Coupled	Uncoupled
Coupled Referential	Direct	Mailbox
Uncoupled	Meeting oriented	Generative communication



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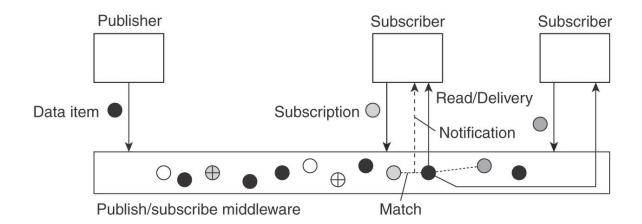
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Jini Case Study

- Coordination system based on Java
 - Clients can discover new services as they become available
 - Example: "intelligent toaster"
 - Distributed event and notification system
- Coordination model
 - Bulletin board model
 - Uses JavaSpaces: a shared dataspace that stores tuples
 - Each tuple points to a Java object



Overall Approach

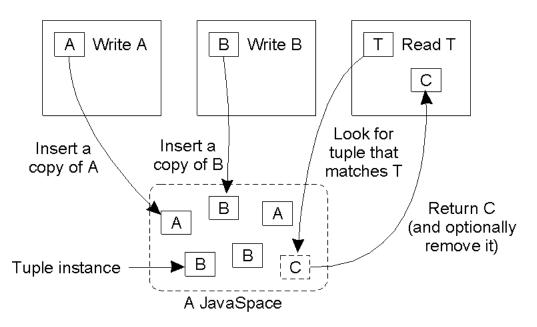


• The principle of exchanging data items between publishers and subscribers.



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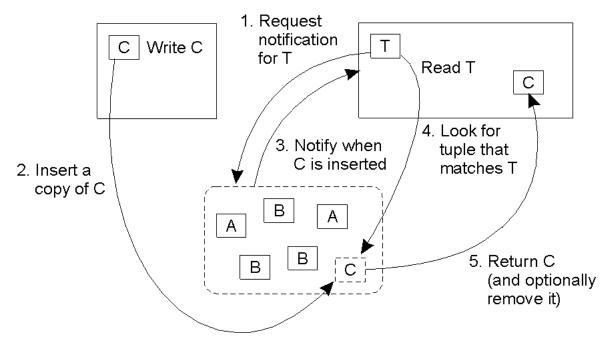
Overview of Jini



The general organization of a JavaSpace in Jini.



Communication Events



Using events in combination with a JavaSpace

Tuple broadcast

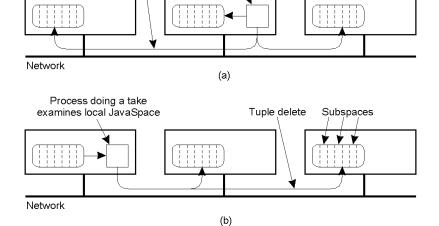


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a write broadcasts



- A JavaSpace can be replicated on all machines. The dotted lines show the partitioning of the JavaSpace into subspaces.
- a) Tuples are broadcast on WRITE
- b) READs are local, but the removing of an instance when calling TAKE must be broadcast



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Processes (2)

Process doing a write inserts tuple into local JavaSpace

Network

(a)

Process doing a read broadcasts template

Template broadcast

Network

(b)

- Unreplicated JavaSpace.
- a) A WRITE is done locally.
- b) A READ or TAKE requires the template tuple to be broadcast in order to find a tuple instance



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