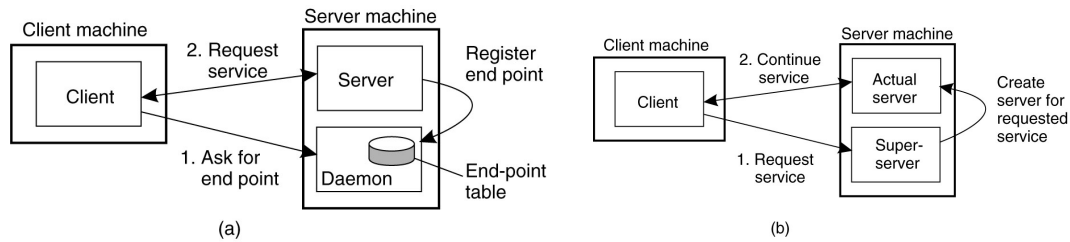


# Server Design Issues



- Server Design
  - Iterative versus concurrent
- How to locate an end-point (port #)?
  - Well known port #
  - Directory service (port mapper in Unix)
  - Super server (inetd in Unix)

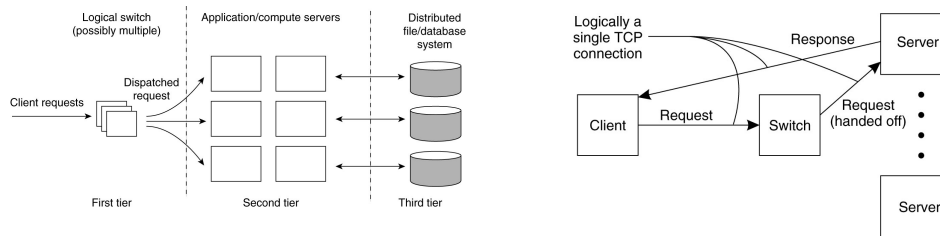


# Stateful or Stateless?

- Stateful server
  - Maintain state of connected clients
  - Sessions in web servers
- Stateless server
  - No state for clients
- Soft state
  - Maintain state for a limited time; discarding state does not impact correctness



# Server Clusters



- Web applications use tiered architecture
  - Each tier may be optionally replicated; uses a dispatcher
  - Use TCP splicing or handoffs

# Code and Process Migration

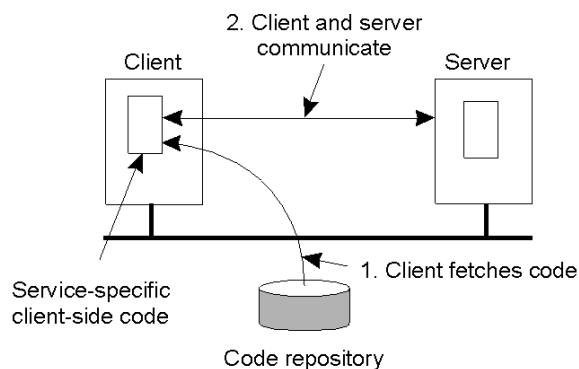
- Motivation
- How does migration occur?
- Resource migration
- Agent-based system
- Details of process migration

# Motivation

- Key reasons: performance and flexibility
- Process migration (aka *strong mobility*)
  - Improved system-wide performance – better utilization of system-wide resources
  - Examples: Condor, DQS
- Code migration (aka *weak mobility*)
  - Shipment of server code to client – filling forms (reduce communication, no need to pre-link stubs with client)
  - Ship parts of client application to server instead of data from server to client (e.g., databases)
  - Improve parallelism – agent-based web searches

# Motivation

- Flexibility
  - Dynamic configuration of distributed system
  - Clients don't need preinstalled software – download on demand



# Migration models

- Process = code seg + resource seg + execution seg
- Weak versus strong mobility
  - Weak => transferred program starts from initial state
- Sender-initiated versus receiver-initiated
- Sender-initiated
  - migration initiated by machine where code resides
    - Client sending a query to database server
      - Client should be pre-registered
- Receiver-initiated
  - Migration initiated by machine that receives code
  - Java applets
  - Receiver can be anonymous

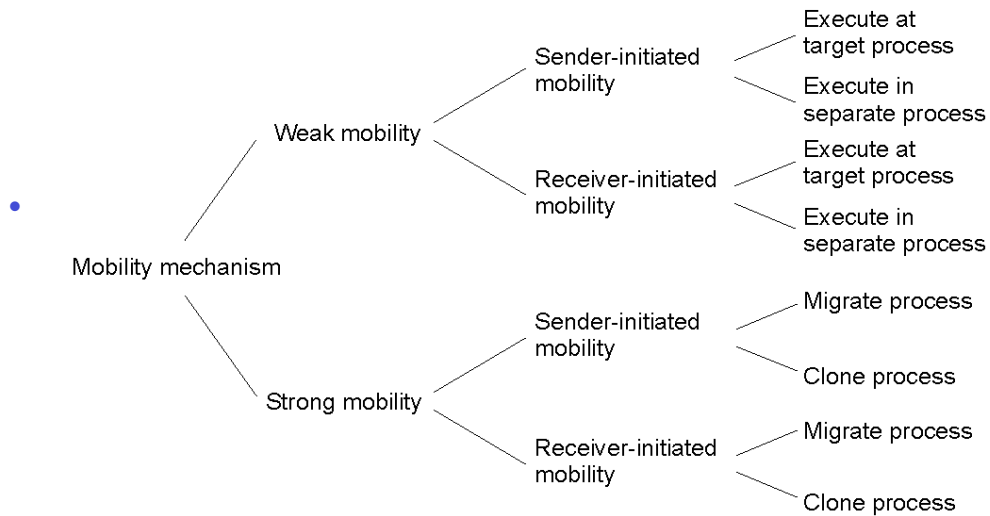


## Who executes migrated entity?

- Code migration:
  - Execute in a separate process
  - [Applets] Execute in target process
- Process migration
  - Remote cloning
  - Migrate the process



# Models for Code Migration



## Do Resources Migrate?

- Depends on resource to process binding
  - By identifier: specific web site, ftp server
  - By value: Java libraries
  - By type: printers, local devices
- Depends on type of “attachments”
  - Unattached to any node: data files
  - Fastened resources (can be moved only at high cost)
    - Database, web sites
  - Fixed resources
    - Local devices, communication end points



# Resource Migration Actions

## Resource-to machine binding

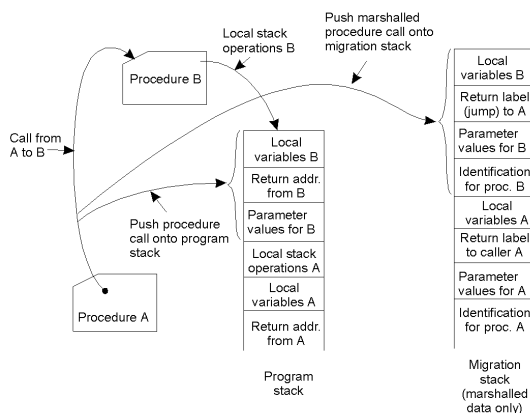
	Unattached	Fastened	Fixed	
Process-to-resource binding	By identifier	MV (or GR)	GR (or MV)	GR
	By value	CP ( or MV, GR)	GR (or CP)	GR
	By type	RB (or GR, CP)	RB (or GR, CP)	RB (or GR)

- Actions to be taken with respect to the references to local resources when migrating code to another machine.
- GR: establish global system-wide reference
- MV: move the resources
- CP: copy the resource
- RB: rebind process to locally available resource



# Migration in Heterogeneous Systems

- Systems can be heterogeneous (different architecture, OS)
  - Support only weak mobility: recompile code, no run time information
  - Strong mobility: recompile code segment, transfer execution segment [migration stack]
  - Virtual machines - interpret source (scripts) or intermediate code [Java]



# Case study: Agents

- Software agents
  - Autonomous process capable of reacting to, and initiating changes in its environment, possibly in collaboration
  - More than a “process” – can act on its own
- Mobile agent
  - Capability to move between machines
  - Needs support for strong mobility
  - Example: D’Agents (aka Agent TCL)
    - Support for heterogeneous systems, uses interpreted languages

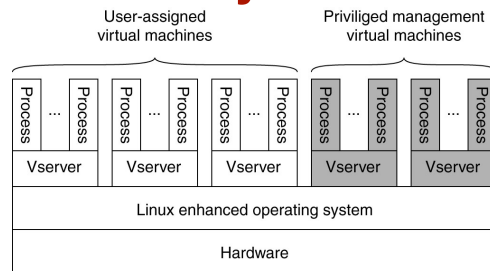


# Case Study: Viruses and Malware

- Viruses and malware are examples of mobile code
  - Malicious code spreads from one machine to another
- Sender-initiated:
  - proactive viruses that look for machines to infect
    - Autonomous code
- Receiver-initiated
  - User (receiver) clicks on infected web URL or opens an infected email attachment



# Case Study: PlanetLab



- Distributed cluster across universities
  - Used for experimental research by students and faculty in networking and distributed systems
- Uses a virtualized architecture
  - Linux Vservers
  - Node manager per machine
  - Obtain a “slice” for an experiment: slice creation service



# Case Study: ISOS

- Internet scale operating system
  - Harness compute cycles of thousands of PCs on the Internet
  - PCs owned by different individuals
  - Donate CPU cycles/storage when not in use (pool resources)
  - Contact coordinator for work
  - Coordinator: partition large parallel app into small tasks
  - Assign compute/storage tasks to PCs
- Examples: [Seti@home](#), P2P backups





# Case study: Condor

- Condor: use idle cycles on workstations in a LAN
- Used to run large batch jobs, long simulations
- Idle machines contact condor for work
- Condor assigns a waiting job
- User returns to workstation => suspend job, migrate
- Flexible job scheduling policies