## Last Class: RPCs and RMI

- Case Study: Sun RPC
- Lightweight RPCs
- Remote Method Invocation (RMI)
  - Design issues



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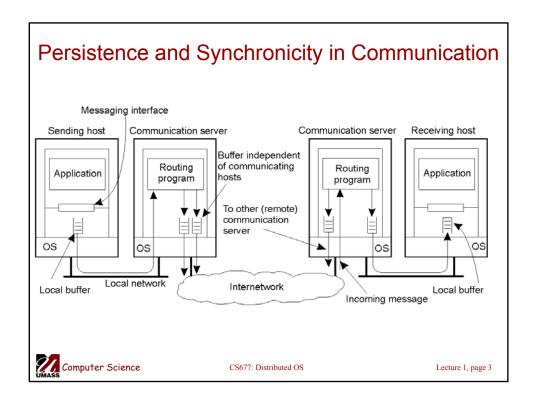
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## Today: Communication Issues

- Message-oriented communication
  - Persistence and synchronicity
- Stream-oriented communication

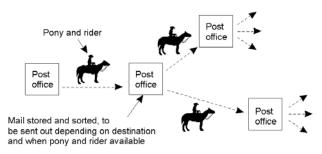


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### **Persistence**

- Persistent communication
  - Messages are stored until (next) receiver is ready
  - Examples: email, pony express



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### **Persistence**

- Transient communication
  - Message is stored only so long as sending/receiving application are executing
  - Discard message if it can't be delivered to next server/receiver
  - Example: transport-level communication services offer transient communication
  - Example: Typical network router discard message if it can't be delivered next router or destination



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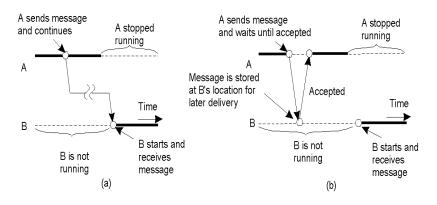
## Synchronicity

- Asynchronous communication
  - Sender continues immediately after it has submitted the message
  - Need a local buffer at the sending host
- Synchronous communication
  - Sender blocks until message is stored in a local buffer at the receiving host or actually delivered to sending
  - Variant: block until receiver processes the message
- Six combinations of persistence and synchronicity



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#### Persistence and Synchronicity Combinations



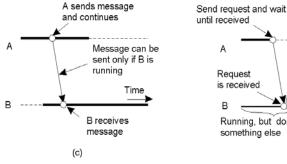
- a) Persistent asynchronous communication (e.g., email)
- b) Persistent synchronous communication

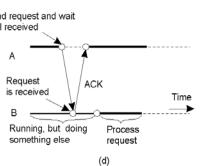


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#### Persistence and Synchronicity Combinations



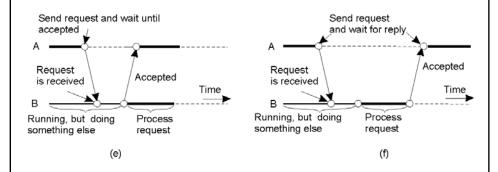


- c) Transient asynchronous communication (e.g., UDP)
- d) Receipt-based transient synchronous communication

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#### Persistence and Synchronicity Combinations



- e) Delivery-based transient synchronous communication at message delivery (e.g., asynchronous RCP)
- f) Response-based transient synchronous communication (RPC)

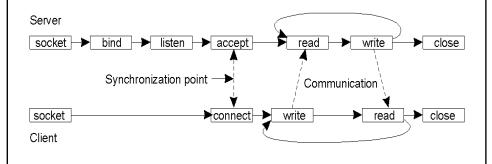


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# Message-oriented Transient Communication

- Many distributed systems built on top of simple message-oriented model
  - Example: Berkeley sockets



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## **Berkeley Socket Primitives**

Primitive	Meaning	
Socket	Create a new communication endpoint	
Bind	Attach a local address to a socket	
Listen	Announce willingness to accept connections	
Accept	Block caller until a connection request arrives	
Connect	Actively attempt to establish a connection	
Send	Send some data over the connection	
Receive	Receive some data over the connection	
Close	Release the connection	



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# Message-Passing Interface (MPI)

- Sockets designed for network communication (e.g., TCP/IP)
  - Support simple send/receive primitives
- Abstraction not suitable for other protocols in clusters of workstations or massively parallel systems
  - Need an interface with more advanced primitives
- Large number of incompatible proprietary libraries and protocols
  - Need for a standard interface
- Message-passing interface (MPI)
  - Hardware independent
  - Designed for parallel applications (uses transient communication)
- Key idea: communication between groups of processes
  - Each endpoint is a (groupID, processID) pair



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#### **MPI** Primitives

Primitive	Meaning	
MPI_bsend	Append outgoing message to a local send buffer	
MPI_send	Send a message and wait until copied to local or remote buffer	
MPI_ssend	Send a message and wait until receipt starts	
MPI_sendrecv	Send a message and wait for reply	
MPI_isend	Pass reference to outgoing message, and continue	
MPI_issend	Pass reference to outgoing message, and wait until receipt starts	
MPI_recv	Receive a message; block if there are none	
MPI_irecv	Check if there is an incoming message, but do not block	



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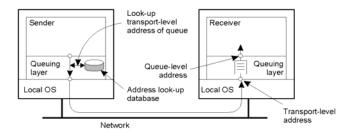
# Message-oriented Persistent Communication

- Message queuing systems
  - Support asynchronous persistent communication
  - Intermediate storage for message while sender/receiver are inactive
  - Example application: email
- Communicate by inserting messages in queues
- Sender is only guaranteed that message will be eventually inserted in recipient's queue
  - No guarantees on when or if the message will be read
  - "Loosely coupled communication"



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## Message-Queuing Model



Primitive	Meaning
Put	Append a message to a specified queue
Get	Block until the specified queue is nonempty, and remove the first message
Poll	Check a specified queue for messages, and remove the first. Never block.
Notify	Install a handler to be called when a message is put into the specified queue.



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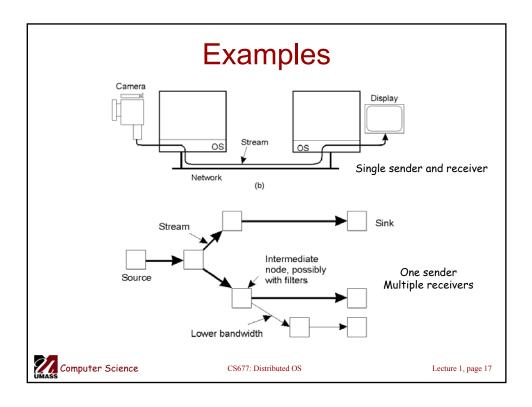
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## **Stream Oriented Communication**

- Message-oriented communication: request-response
  - When communication occurs and speed do not affect correctness
- Timing is crucial in certain forms of communication
  - Examples: audio and video ("continuous media")
  - 30 frames/s video => receive and display a frame every 33ms
- Characteristics
  - Isochronous communication
    - Data transfers have a maximum bound on end-end delay and jitter
  - Push mode: no explicit requests for individual data units beyond the first "play" request

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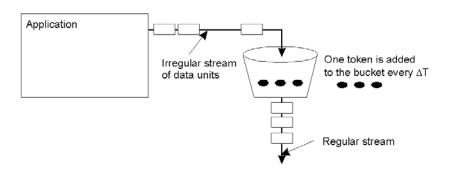


# Quality of Service (QoS)

- Time-dependent and other requirements are specified as *quality of service (QoS)* 
  - Requirements/desired guarantees from the underlying systems
  - Application specifies workload and requests a certain service quality
  - Contract between the application and the system

Characteristics of the Input	Service Required
•maximum data unit size (bytes)	•Loss sensitivity (bytes)
•Token bucket rate (bytes/sec)	•Loss interval (μsec)
•Toke bucket size (bytes)	<ul> <li>Burst loss sensitivity (data units)</li> </ul>
Maximum transmission rate	<ul> <li>Minimum delay noticed (μsec)</li> </ul>
(bytes/sec)	<ul> <li>Maximum delay variation (μsec)</li> </ul>
	•Quality of guarantee

# Specifying QoS: Token bucket



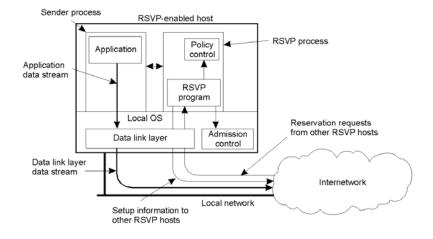
- The principle of a token bucket algorithm
  - Parameters (rate r, burst b)
  - Rate is the average rate, burst is the maximum number of packets that can arrive simultaneously



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# Setting Up a Stream: RSVP





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