Last Class: Distributed Systems and RPCs

- Servers export procedures for some set of clients to call
- To use the server, the client does a procedure call
- OS manages the communication



CS377: Operating Systems

Lecture 22, page 1

Today: Distributed File Systems

- One of the most common uses of distributed systems
- Basic idea:
 - Given a set of disks attached to different nodes.
 - share disks between nodes as if all the disks were attached to every node.
- Examples:
 - Edlab: One server node with all the disks, and a bunch of diskless workstations on a LAN.
 - AppleShare: Every node is both a server with a disk and a client.



CS377: Operating Systems

Distributed File Systems: Issues

- Naming and Transparency
- Remote file access
- Caching
- Server with state or without
- Replication



CS377: Operating Systems

Lecture 22, page 3

Naming and Transparency

- Issues
 - How are files named?
 - Do file names reveal their location?
 - Do file names change if the file moves?
 - Do file names change if the user moves?
- **Location transparency:** the name of the file does not reveal the physical storage location.
- **Location independence:** The name of the file need not change if the file's storage location changes.
- Most naming schemes used in practice do not have location independence, but many have location transparency.



CS377: Operating Systems

Naming Strategies: Absolute Names

- Absolute names: <machine name: path name>
- Examples: AppleShare, Win NT
- Advantages:
 - Finding a fully specified file name is simple.
 - It is easy to add and delete new names.
 - No global state.
 - Scales easily.

Disadvantages:

- User must know the complete name and is aware of which files are local and which are remote.
- File is location dependent, and thus cannot move.
- Makes sharing harder.
- Not fault tolerant.



CS377: Operating Systems

Lecture 22, page 5

Naming Strategies: Mount Points

- Mount Points (NFS Sun's Network File System)
 - Each host has a set of local names for remote locations.
 - Each host has a mount table (/etc/fstab) that specifies <remote path name @ machine name> and a <local path name>.
 - At boot time, the local name is bound to the remote name.
 - Users then refer to the <local path name> as if it were local, and the NFS takes care of the mapping
- Advantages: location transparent, remote name can change across reboots
- **Disadvantages:** single unified strategy hard to maintain, same file can have different names



CS377: Operating Systems

NFS: Example

Partial contents of /etc/fstab for Edlab machines:

/usr1/mail@elux3.cs.umass.edu:/var/spool/mail

/users/users1@elsrv1:/users/users1

/users/users2@elsrv1:/users/users2

/users/users3@elsrv2:/users/users3

/users/users4@elsrv2:/users/users4

/courses/cs300@elsrv3:/courses/cs300

/rcf/mipsel/4.2/share@elsrv1:/exp/rcf/share

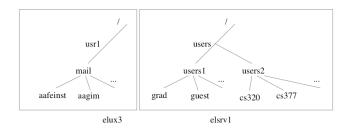
/rcf/common@elsrv1:/exp/rcf/common

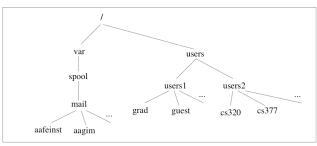


CS377: Operating Systems

Lecture 22, page 7

NFS: Example





deluxa



CS377: Operating Systems

Naming Strategies: Global Name Space

- Single name space: CMU's Andrew and Berkeley's Sprite
 - No matter which node you are on, the file names are the same.
 - Set of workstation clients, and a set of dedicated file server machines.
 - When a client starts up, it gets its file name structure from a server.
 - As users access files, the server sends copies to the workstation and the workstation caches the files



CS377: Operating Systems

Lecture 22, page 9

Global Name Space

Advantages:

- Naming is consistent and easy to keep consistent.
- The global name space insures all the files are the same regardless of where you login.
- Since names are bound late, moving them is easier.

Disadvantages:

- It is difficult for the OS to keep file contents consistent due to caching.
- Global name space may limit flexibility.
- Performance problems.



Remote File Access and Caching

Once the user specifies a remote file, the OS can do the access either

- 1. remotely, on the server machine and then return the results using RPC (called *remote service*), or
- 2. can transfer the file (or part of the file) to the requesting host, and perform local accesses (called *caching*)

Caching Issues:

- Where and when are file blocks cached?
- When are modifications propagated back to the remote file?
- What happens if multiple clients cache the same file?



CS377: Operating Systems

Lecture 22, page 11

Remote File Access and Caching

Location

- Local disk
 - Advantages:
 - Access time reduced.
 - Safer if node fails.
 - Disadvantages:
 - Difficult to keep local copy consistent with remote copy.
 - Slower than just keeping it in local memory.
 - Requires client to have a disk.



CS377: Operating Systems

Remote File Access and Caching

Location

- Local memory
 - Advantages: Quick access time.
 - Disadvantages:
 - Difficult to keep local copy consistent with remote copy.
 - Does not tolerate node failure well.
 - Limited cache size.
 - Works with diskless workstations.



CS377: Operating Systems

Lecture 22, page 13

Cache Update Policies

When to write local changes to the server has a central role in determining distributed file system performance.

- Write through: yields the most reliable results since every write hits the remote disk before the process continues, but it has the poorest performance.
 - Caching with write through is equivalent to using remote service for all writes, and exploits caching only for reads.
- **Write back:** yields the quickest response time since the write need only hit cache before the process continues.
 - It reduces network traffic and the number of writes to the disk for repeated writes to the same disk block, since only one of the writes will go across the network.
 - If a user machine crashes, the unwritten data is lost.
 - Write-back when file is closed, a block is evicted from cache, or every 30sec.



CS377: Operating Systems

Cache Consistency

- Client-initiated consistency: Client contacts the server and asks if its copy is consistent with the server's copy.
 - Can check every access.
 - Can check at a given interval.
 - Can check only upon opening a file.
- **Server-initiated consistency:** Server detects potential conflicts and invalidates caches
 - Server needs to know:
 - which clients have cached which parts of which files.
 - which clients are readers and which are writers.



CS377: Operating Systems

Lecture 22, page 15

Server State and Replication

- Stateful versus stateless server
 - Web analogy
 - Tradeoff between performance and tolerence to crash faults
- Replication
 - Server data is replicated across machines
 - Need to ensure consistency of files when a file is updated on one server



CS377: Operating Systems

Case Study: Sun's Network File System

- NFS is the standard for distributed UNIX file access.
- NFS is designed to run on LANs (but works on WANs)
- Nodes are both servers and clients.
- Servers have no state (NFS v3 only; NFS v4 is stateful)
- Uses a mount protocol to make a global name local
 - /etc/exports lists the local names the server is willing to export.
 - /etc/fstab lists the global names that the local nodes import. A corresponding global name must be in /etc/exports on the server.



CS377: Operating Systems

Lecture 22, page 17

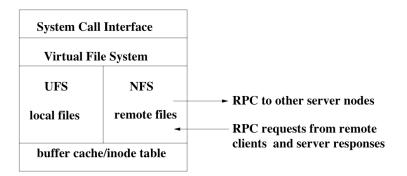
NFS Implementation

- NFS defines a set of RPC operations for remote file access:
 - 1. directory search, reading directory entries
 - 2. manipulating links and directories
 - 3. accessing file attributes
 - 4. reading/writing files
- Does not rely on node homogeneity heterogeneous nodes must simply support the NFS mount and remote access protocols using RPC.
- Users may need to know different names depending upon the node to which they logon.



CS377: Operating Systems

NFS Implementation





CS377: Operating Systems

Lecture 22, page 19

NFS Implementation

- NFS defines new layers in the Unix file system
- The virtual file system provides a standard interface, using vnodes as file handles. A vnode describes either a local file or a remote file.
- The ``buffer cache' caches remote file blocks and attributes.
- On an *open*, the client asks the server whether its cached blocks are up to date.
- Once a file is open, different clients can write to it and get inconsistent data.
- Modified data is flushed back to the server every 30s.
- What file contents do new clients see?
 - Effects of last flush. Writers might have made changes but not updated remote file yet.
- What file contents do existing clients see?
 - For cached blocks, they see out of date info. For new blocks, same as new client



CS377: Operating Systems

Summary

- Naming
 - Desire name independence, but it is difficult to attain
 - Location dependent names are most prevalent
- Speed up remote file access with caching
- Need to write changes back to disk



CS377: Operating Systems